

Type Inference for Units of Measure

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Trends in Functional Programming

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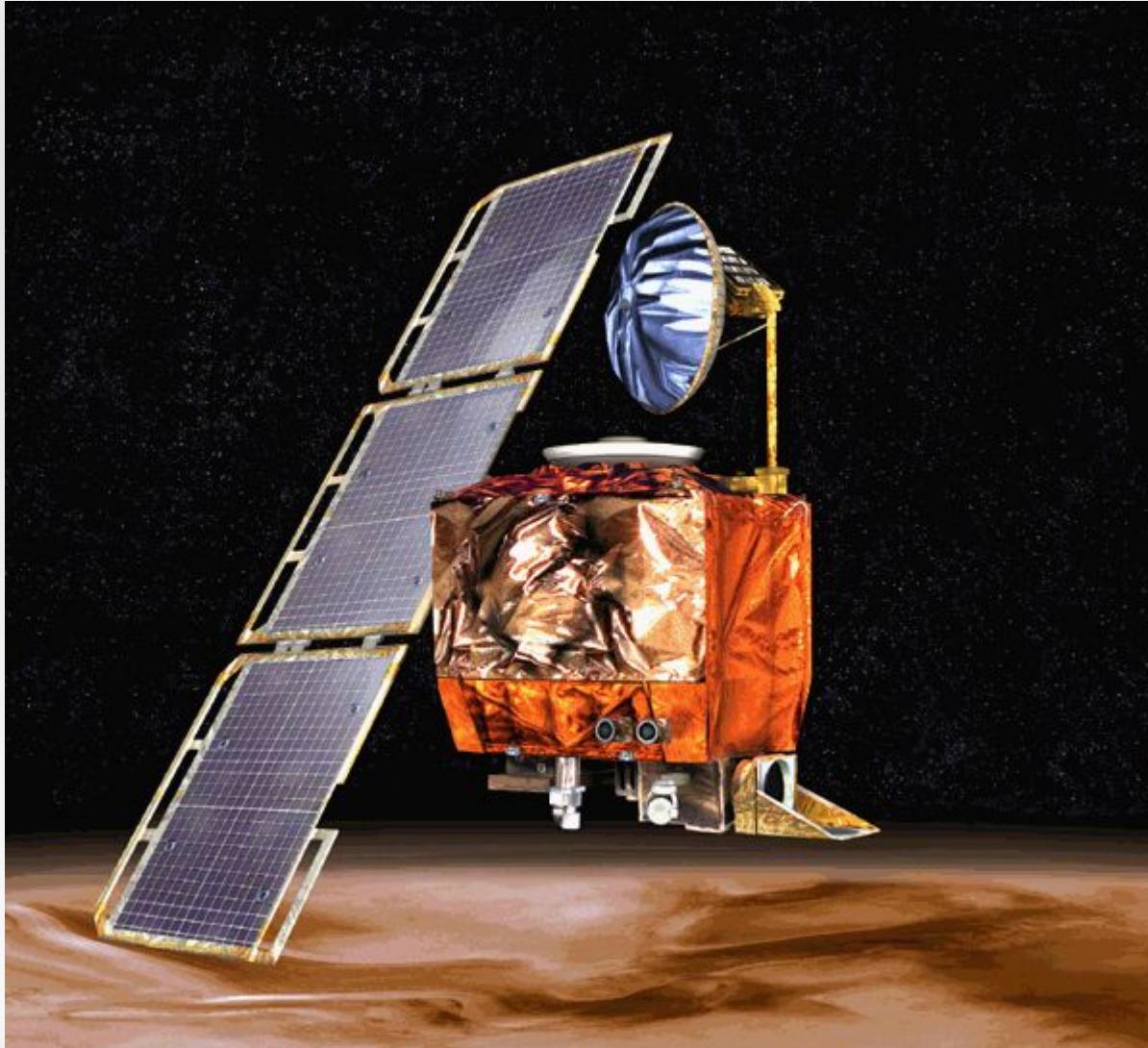
What are units of measure?

- A **dimension** is a physical quantity
 - length
 - time
 - mass
- A **unit** is a standard measure of quantity
 - metres
 - feet
 - seconds

What are units of measure?

- Arithmetic only works if units are compatible:
 - $10 \text{ m} + 5 \text{ m} = 15 \text{ m}$
 - $120 \text{ m} / 60 \text{ s} = 2 \text{ ms}^{-1}$
 - $6 \text{ m} - 3 \text{ s} = ???$
- Can we enforce unit compatibility with types?

Why?



NASA/JPL

The plan

- Algebraic structure of units
- Units of measure in F#
- Type inference going wrong
- Type inference done differently
- Future speculations

Algebraic structure

- Base units
 - metres (m)
 - seconds (s)
 - ...
- Derived units
 - square metres (m^2)
 - metres per second (ms^{-1})
 - ...

Algebraic structure

- We have:
 - Multiplication: e.g. $m^2 = m \cdot m$
 - Dimensionless quantities: 1
 - Inverses: e.g. s^{-1}
- Subject to:
 - $d \cdot (e \cdot f) = (d \cdot e) \cdot f$ (associativity)
 - $1 \cdot d = d = d \cdot 1$ (identity)
 - $d \cdot d^{-1} = 1$ (inverse)
 - $d \cdot e = e \cdot d$ (commutativity)

Algebraic structure

- Units form an abelian group
- Specifically, the **free** abelian group generated by the base units
- No fractional powers...
- ...but we probably don't need them (?)

Units of measure in F#

- Andrew Kennedy pioneered work on units of measure with **polymorphism**
- He introduced them in F#
- I'm following his design

Units of measure in F#

```
type [] m;
type [] s;
let vel    = 2.0<m/s>;
let accel = 3.8<m/s^2>;
let distance t =
    vel * t + accel * t * t;
val distance : float<s> → float<m>
```

Type inference is possible

- Free abelian group unification
 - has most general unifiers
 - is decidable
- We can infer types with Damas and Milner's Algorithm W

Type inference is tricksy

```
> fun x ->  
-     (div x 5<m>, div x 2<s>);;  
  
val it : int<'u> -> int<'u/m> * int<'u/s>
```

Type inference is tricksy

```
> fun x -> let f = div x in  
-     (f 5<m>, f 2<s>);;
```

Type inference is tricksy

```
> fun x -> let f = div x in  
-      (f 5<m>, f 2<s>);;  
----- ^^^^
```

error FS0001: Type mismatch. Expecting a

int<m>

but given a

int<s>

The unit of measure 'm' does not match the unit of measure 's'

Type inference is tricksy

- F# doesn't always infer principal types
- Let-generalisation is syntactic:
- does **a** occur free in the typing environment?
- This doesn't respect group equivalence:
- e.g. $a \cdot a^{-1} \equiv 1$ but **a** only occurs on one side

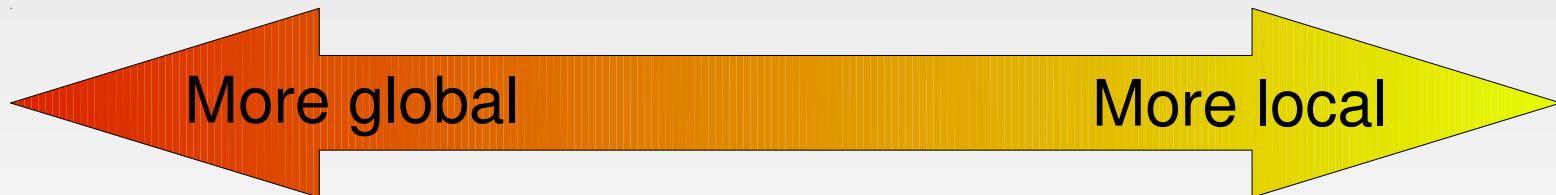
Type inference going wrong

```
fun x -> let f = div x in (f 5<m>, f 2<s>) : ?  
x : t          ⊢ let f = div x in (f 5<m>, f 2<s>) : ?  
x : t          ⊢ div x : ?  
x : t          ⊢ div : int<a b> → int<a> → int<b>  
x : t          ⊢ div x : int<a> → int<b> (if t = int<a b>)  
x : int<a b>    ⊢ div x : int<a> → int<b>  
x : int<a b>, f : int<a> → int<b> ⊢ (f 5<m>, f 2<s>) : ?  
x : int<a b>, f : int<a> → int<b> ⊢ f 5<m> : int<b> (if a = m)  
x : int<m b>, f : int<m> → int<b> ⊢ f 2<s> : int<b> (if m = s)  ✘
```

Type inference done differently

- Types go in the context
- Ordered by dependency

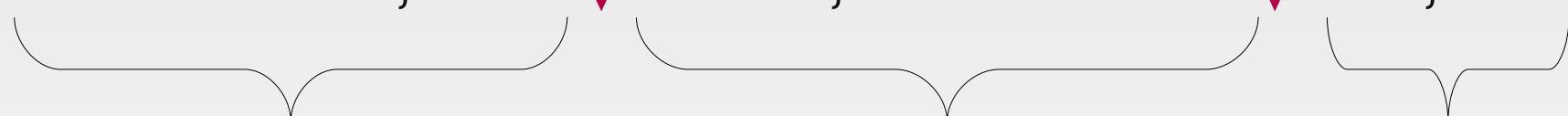
`a := int<m>, ?b, x : b, c := a → b, ?d, ...`



Type inference done differently

- Context divided into 'localities'
- Mark generalisation points for let-expressions

a := int<m>, ?**b** ♦ **x** : **b**, **c** := **a** → **b** ♦ ?**d**, ...



Type inference done differently

- Type variables only moved when necessary
- Most general unifier is a more precise notion
- Generalisation is easy: collect variables from the current locality

Type inference example

```
fun x -> let f = div x in (f 5<m>, f 2<s>) : ?
```

Type inference example

```
fun x -> let f = div x in (f 5<m>, f 2<s>) : ?
```

```
?t, x : t      ⊢ let f = div x  
                in (f 5<m>, f 2<s>) : ?
```

Type inference example

fun x -> let f = div x in (f 5<m>, f 2<s>) : ?

?t, x : t ♦

$\vdash \text{div } x : ?$

?t, x : t ♦ ?a, ?b

$\vdash \text{div} : \text{int}<\text{a b}> \rightarrow \text{int}<\text{a}> \rightarrow \text{int}<\text{b}>$

?t, x : t ♦ ?a, ?b

$\vdash \text{div } x : \text{int}<\text{a}> \rightarrow \text{int}<\text{b}> \text{ (if } t = \text{int}<\text{ab}>\text{)}$

?t, x : t ♦ ?a, ?b, ?c

$\vdash \text{div } x : \text{int}<\text{a}> \rightarrow \text{int}<\text{b}> \text{ (if } t = \text{int}<\text{c}>, c = \text{a b}\text{)}$

?c, x : int<c> ♦ ?a, ?b

$\vdash \text{div } x : \text{int}<\text{a}> \rightarrow \text{int}<\text{b}> \text{ (if } c = \text{a b}\text{)}$

?c, x : int<c> ♦ ?a, b := c a⁻¹

$\vdash \text{div } x : \text{int}<\text{a}> \rightarrow \text{int}<\text{b}>$

?c, x : int<c>

$\vdash f : \forall a. \text{int}<\text{a}> \rightarrow \text{int}<\text{c a}^{-1}>$

?c, x : int<c>, f : ...

$\vdash (f 5<\text{m}>, f 2<\text{s}>) : \text{int}<\text{c m}^{-1}> \times \text{int}<\text{c s}^{-1}>$

$\vdash \dots : \forall c. \text{int}<\text{c}> \rightarrow \text{int}<\text{c m}^{-1}> \times \text{int}<\text{c s}^{-1}>$

Where do we go from here?

- Another free abelian group: the integers
- I'm using this approach to type inference for:
 - Numeric inequalities
 - Local constraints (GADTs)
 - Higher-rank types
 - Lexically-scoped type variables

NEW CUYAMA

Population

562

Ft. above sea level

2150

Established

1951

TOTAL 4663

References

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- George Kuan and David MacQueen
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